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マス·フォア·インダストリ研究 No.21 Fiber Topology Meets Applications

Organizer:

Daisuke Sakurai Shigeo Takahashi Naoki Hamada Osamu Saeki Hamish Carr Takahiro Yamamoto Institute of Mathematics for Industry Kyushu University

About the Mathematics for Industry Research

The Mathematics for Industry Research was founded on the occasion of the certification of the Institute of Mathematics for Industry (IMI), established in April 2011, as a MEXT Joint Usage/Research Center – the Joint Research Center for Advanced and Fundamental Mathematics for Industry – by the Ministry of Education, Culture, Sports, Science and Technology (MEXT) in April 2013. This series publishes mainly proceedings of workshops and conferences on Mathematics for Industry (MfI). Each volume includes surveys and reviews of MfI from new viewpoints as well as up-to-date research studies to support the development of MfI.

October 2018 Osamu Saeki Director Institute of Mathematics for Industry

Fiber Topology Meets Applications

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Fiber Topology Meets Applications

Organizer: Daisuke Sakurai Shigeo Takahashi Naoki Hamada Osamu Saeki Hamish Carr Takahiro Yamamoto

Preface

In differential topology, the study of the topology of fibers, known as fiber topology, extends the Morse theory, of functions, to that of maps. This establishes a thorough exploration of topological transitions of inverse images. It has been proven powerful for understanding today's data whose size and complexity are overwhelming.

In recent years, fiber topology has been actively pushed forward from mathematical studies towards computation. While the computational approach for the Morse-case was studied separately, the two research domains started communication after some years. For more general cases, theoretical insights have been transferred into numerical algorithms through direct collaboration with computer scientists. This brought a bi-directional impacts where (i) mathematics boosts the progress in computation and (ii) computer science also helps mathematics. Simultaneously, fiber topology has been applied for decision making with multiple optimization problems.

The organizing board is now confident in adapting the achieved theories and techniques for industrial and inter-governmental decision making. We have conducted discussions with world-leading research groups to identify mutual interests that would advance both the society and techniques. With the proposed meeting, we finish the initial goal setting phase, and move on to the mission of providing solutions to application problems.

The meeting aids the interaction among mathematics, computation, and applications. For this, participants gather from broad spectra of research disciplines, including not only mathematics and computer science, but also operational research, sciences for environmental and semiconductors.

The operational research, led by Naoki Hamada (KLab), Daisuke Sakurai (Kyushu University Pan-Omics Data-Driven Research Innovation Center) and Takahiro Yamamoto (Tokyo Gakugei University) aims at designing benchmark problems for multioptimization problems utilizing knowledge extracted from fiber topology.

Environmental scientists Hiroshi Yamashita and Bastian Kern (German Aerospace Center, Germany) have some history of joint-research with Sakurai while he was at Zuse Institute Berlin (Germany). The current research project focuses on proposing a new tax incentive for balancing the environmental costs against the operational cost in the aviation industry. Fiber topology is employed for understanding the high-dimensional space spanned by different kinds of subcosts.

Finally, the application in semiconductor research employs fiber topology for understanding the potential fields of atomic configurations. For simulation outputs done with Japan's next-generation supercomputer Fugaku, we extract the landscape of the high-dimensional potential fields. Fiber topology gives a concise and parameter-free representation of the landscape, which is also robust against numerical errors.

All of the applications pose a challenge to the current state-of-the-art analysis of fiber topology. Especially, the high spatial dimensionality in the data is something that requires an update to the numerical methods. The current, rather loose, mathematics found in the computer science literature meets a more rigorous treatment by Takahiro Yamamoto (Tokyo Gakugei University).

Theorists not only provide solutions, but get feedback about desirable expansions. We discuss how fiber topology should let us extract information about the energy potential of atomic configurations, robustness of the assembly schedule against errors, and the impact of the inter-governmental decisions towards the aviation industry and air pollution.

While our progress in computation has seen a success, there are theoretical questions left, such as how the mathematical results translate to and from numerical representations.

What directions in mathematical theories benefit numerical ones, and vice versa, is another question.

The forum was co-sponsored by IMI and Pan-Omics Data-Driven Research Innovation Center. It was also supported by Kyushu University Kasseika project (令和2年度九州大学活性化制度).

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Fiber Topology Meets Applications

日時: 2021年1月6日(水)14:30~2021年1月8日(金)16:45

場所: 九州大学 伊都キャンパス ウエスト1号館 D棟4階 IMIオーディトリアム(W1-D-413)

※Zoomミーティングを使用したハイブリッド型で開催されます https://www.imi.kyushu-u.ac.jp/kyodo-riyo/research_meetings/view/13

▶ 下記URLより参加登録をお願いいたします https://forms.gle/EBEQP7VjivtienPEA

1月6日(水)

14:30-15:15

Title : Fiber Topology Meets Applications: Where We are Heading Speaker : Daisuke Sakurai (Kyushu University)

15:15-16:00

Title : Why Topology is Necessary at Exascale - And Why it's Not Easy Speaker : Hamish Carr (University of Leeds)

16:00-16:45 Title : Modeling semiconductor epitaxy for next generation power device application Speaker : Yoshihiro Kangawa (Kyushu University)

1月7日(木)

14:30-15:15

Title : Data Parallel Hypersweeps for in Situ Topological Analysis Speakers : Petar Hristov (University of Leeds), Gunther Weber (Lawrence Berkeley National Laboratory and University of California, Davis), Hamish Carr (University of Leeds), Oliver Rübel (Lawrence Berkeley National Laboratory) and James Ahrens (Los Alamos National Laboratory)

15:15-16:00

Title : Design of Multimodal Test Problems in Multiobjective Optimization Using Fiber Topology

Speakers : Reiya Hagiwara (Kyushu University), Naoki Hamada (KLab Inc.), Takahiro Yamamoto (Tokyo Gakugei University), Daisuke Sakurai (Kyushu University)

16:00-16:45 Title : Eco-efficient Flight Trajectory Exploration by Using the Chemistry-climate Model EMAC Speakers : Hiroshi Yamashita and Bastian Kern (German Aerospace Center)

1月8日(金)

14:30-15:15 Title : Elimination of B2 singularities Speaker : Takahiro Yamamoto (Tokyo Gakugei University)

15:15-16:00 Simplifying Indefinite Fibrations on 4-manifolds Speaker : Osamu Saeki (Kyushu University)

16:00-16:30 Title : An Efficient Triangulation for Extruded Spatiotemporal Prism Meshes Speakers : Akito Fujii, Kenji Ono and Daisuke Sakurai (Kyushu University)

16:30-16:45 Concluding Remarks Daisuke Sakurai (Kyushu University)

Fiber Topology Meets Applications: Where We Are Heading

Daisuke Sakurai

Kyushu Univ 2020 January 6

Thanks For Joining!

- We Are From
 - Math (Singularity Theory)
 - · CompSci (Visualization)
 - Application Areas (Materials, Gaming, Atmosphere)

Thanks For Joining!

- · Mutual Interest in
 - · Topology of Fibers, eg
 - · Configuration of
 - · Local Extrema
 - Energy Landscapes

Thanks For Joining!

3

- Share Information
- · Co-design Future Research



Fiber = Inverse Image

Space-Time $f: \mathbb{R}^n \to \mathbb{R}^m$

Measurements

 $f: \mathbb{R}^2 \to \mathbb{R}^1$ [gisgeography.com]

 $f: \mathbb{R}^3 \to \mathbb{R}^1$





 $f: \mathbb{R}^3 \to \mathbb{R}^m$ (Multi-Fields)

Fibers: $f^{-1}(f_1, f_2)$











Math Development (O Saeki)

- Understand Fibers
- → Reliable Algorithms



Classification & Configuration of Fibers, Morphs of Maps, ...









Why Topology Is Needed ... And Why It's Not Easy

Hamish Carr

UNIVERSITY OF LEEDS

Topology

<section-header><section-header><list-item><list-item><list-item><table-container><list-item><table-container>



Optic Throughput		
• 20KB	1s	
• 1MB	50s ~ 1 m	
• 1GB	50,000s ~ 13.9h	
• 1TB	50,000,000s ~ 1.5y	
• 1PB	$50,000,000,000s \sim 1,500y$	
• 1EB	$50,000,000,000,000s \sim 1,500,000y$	
Торо	UNIVERSITY OF LEEDS	

A Visualisation Challenge

- 1 EB of data feasible (10^{18})
- 1 GB limit to comprehension (10⁹)
- 9 orders of magnitude (10⁻⁹) data reduction
 - More if possible
- Visualisation turns into analysis
 - Domain-specific or domain-independent
 - Almost always based on mathematics

Topology



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Combinatorial Reduction

- Our analysis uses continuous mathematics
- Our computers use discrete mathematics
- We rely on mathematical simplifications:
 - Meshes, particles, &c. to discretise the mathematics
 - Numerics to compute the fields
 - Filtrations to enforce sequence

Topology

• Most commonly, filtration by function value

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Simplification / Branch Decomposition

- Large trees have millions of edges
- So we prune edges in sequence
- With importance measure
- Results in a branch decomposition

Topology

Sample Terrain and Tree Region Colours Match Edge Colours

Area

Height

-19-

Serial vs. Parallel

- Filtrations usually process data in sequential sorted order
- We add edges to the tree one at a time
- Simplification is one edge at a time
- Proofs use linear induction
- So the big question is:
 - How do we scale up to parallel systems?
 - Petar will discuss this tomorrow

Topology

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Acknowledgements

- NSERC
- University of British Columbia
- SFI
- University College Dublin
- EPSRC EP/J013072/1
- University of Leeds
- DoE/NNSA ECP Alpine 17-SC-20-SC LBNL Subcontract 7452335

Topology






























Data Parallel Hypersweeps for In Situ Computation

Petar Hristov Gunther Weber Hamish Carr Oliver Ruebel James Ahrens

Overview

1. In Situ Computation

2. Parallel Contour Tree Analysis

3. In Situ Contour Tree Pipeline



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Evolution of Algorithms at Scale

– Mathematical Model

Single Core Algorithm

Shared Memory Parallel Algorithm

*Hybrid Distributed and Shared Memory Algorithm

- Reeb Graph (1940s).
- Sweep and Merge Algorithm (2000s).
- Parallel Peak Pruning (2010s).
- Distributed Parallel Peak Pruning (2020s).



Parallel Algorithm Principles

- Use local independent operations.
- Rely on parallel primitives test, scan, sort.
- Avoid synchronisation and global invariants.
- Avoid data transfers between devices.



Parallel Peak Pruning

- Instead of incremental computation
- Compute monotone paths in parallel
 - Essentially, up the gradient lines
- Use them to identify peaks
- Strip out all peaks at once
- Repeat recursively













Contour Tree Phase

- Essentially the same as the serial
- But batched, alternating upper / lower leaves
- Vertex deletion is lazy: we set a flag
- Then use path compression to remove them
- This allows *parallel* transfer of leaves







Visualizing a Laser Simulation



WarpX Simulation: Data Type = Scalar Field Data Size = 7GB Data Dimensions = 6791 x 371 x 371

Problems with multiple Isosurfaces:

- Manual selection of isovalues. Occlusion of connected components.
- Noise and unwanted features.

Multiple isosurfaces (Cluttered) Top ten individual contours (Clear)



- Contour Tree analysis:
- Compute the Contour Tree 1
- Simplify the Contour Tree 2

3. Extract contours using the Contour Tree

Advantages of contour tree analysis

Limited user interaction.

- Automatic removal of noise.
- Supports multiple importance metrics.



Performance & **Evaluation** Hypersweep seconds 0.005 Branch Decomp Compute Tree seconds Contour Tree Ratio Ratio Supernodes 288.807 HS/CT BD/CT Dataset Dimensions seconds 6791x371x371 1024x1024x750 0.055 7.299 0.01% 0.01% 2.20% WarpX_E_x 317.191 Spathorhynchus 44,554,912 330.926 0.459 1024x1024x795 958x646x1088 55,778,125 89,117,386 268.833 352.491 0.589 0.21% 3.30% 3.92% Kingsnake 8.887 13.841 Pawpawsaurus



Acknowledgments

This research was supported by:

- Exascale Computing Project (17-SC-20-SC), a collaborative effort of the U.S. Department of Energy Office of Science and the National Nuclear Security Administration under Contract No. DE-AC02-05CH11231 (Lawrence Berkeley National Laboratory), Award Number 14-017566 (Los Alamos National Laboratory), and Subcontract 7452335 (University of Leeds).
- School of Computing at the University of Leeds (Scholarship to P. Hristov)
- Resources of the National Energy Research Scientific Computing Center (NERSC), a U.S. Department of Energy Office of Science User Facility operated under Contract No. DE-AC02-05CH11231.

We thank Jean-Luc Vay and Maxence Thevenet for making the WarpX dataset available.































Modality: three definitions

- Definition by local Pareto optima [folklore]
 A problem is multimodal if it has a local Pareto optimum that is not globally Pareto optimal.
- Definition by component modality [Huband+ 2006]
 A problem is *multimodal* if an objective function has multiple local optima.
- Definition by local Pareto set [Kerschke+ 2016]
 A problem is multimodal if its local Pareto set consists of multiple connected components.

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DLR.de • Chart 4 > Eco-efficient Flight Trajectories > Yamashita, Kern • IMI Wo	rkshop > 06.01.2021
National and	international networking
Clients and partners:	Governments and ministries, agencies and organisations, industry and business, science and research
Worldwide 🔫	
Europe	
Germany	HELMHOLTZ Implementation ICEMEINSCHAFT Bundesministerium Bundesministerium Bundesministerium Bundesministerium Bundesministerium Bundesministerium Bundesministerium Bundesministerium Bundesministerium
	Deutsches Zentrum DLR für Luft- und Raumfahrt
ULR	







DLR.de • Chart 8 > Eco-efficient Flight Trajectories > Yamashita, Kern • IMI Workshop > 06.01.2021

The DLR Institute of Atmospheric Physics



DLR

Physics and chemistry of the global atmosphere: 0-120 km altitude

Socially relevant issues related to the atmosphere in aviation, space travel, transport and energy

Climate protection, mobility of the future, digitalization & artificial intelligence, energy system transformation

Both basic and application-oriented questions

Broad spectrum of methods

Internationally competitive and in some areas internationally leading

Competent contact for DLR, society, industry and politics



DLR.de • Charl 9 > Eco-efficient Flight Trajectories > Yamashita, Kern • IMI Workshop > 06.01.2021
The institute at a glance



Founded 1.7.1962 (1924) End of 2019 150 employees (51f, 99m) thereof ~ 37 PhD students 18 Lectureships/professorships at 9 universities/colleges

Overall budget 2019: 18,8 M€ (~ 2256 M¥) Basic funding : 13,2 M€ (48% Aerospace, 39% Aviation, 11% Traffic, 2% Energy) Third-party funds : 5,6 M€ (ESA, EU/ERC, BMBF, BMWi, DFG, HGF, Airbus,...)

DLR.de • Chart 10 > Eco-efficient Flight Trajectories > Yamashita, Kem • IMI Workshop > 06.01.2021 Organization						
	Logistics Sonja Koch Stefanie Zähnle Controlling	Institute of Atmospheric Physics Prof. Markus Rapp Director	Quality Management Dr. Arthur Schady			
Earth System Modelling Prof. Robert Sausen YIG MACClim JProf. Hella Garny	ESM – Evaluation and Analysis Prof. Veronika Eyring	Atmospheric Trace Species Dr. Anke Roiger	Cloud Physics Prof. Christiane Voigt Passive Cloud- Remote Sensing Dr. Luca Bugliaro	Transport Meteorology Dr. Thomas Gerz	Lidar Dr. Andreas Fix Lidar/Radar-Synergy Dr. Silke Groß	
a) New appointment as of 01.07.2021; W2 appointment procedure with LMU ongoing close cooperation with						
DLR		Met self	and the second sec			



https://www.dkrz.de/about/media/galerie/Media-DKRZ/hlre-3











Contents

- Aviation and climate impact
- Climate-optimized routing
- Research objectives and methodologies
- EMAC/AirTraf model components
- 1-day air traffic simulations over the North Atlantic with different aircraft routing options
- Multi-objective optimization in EMAC/AirTraf
- Summary research topics for further collaborations








Climate-optimized routing

- "Climate cost function (CCF) identifies climate sensitive regions for emissions (CO₂, H₂O, ozone, Methane, contrails) and estimates climate impacts
- Climate-optimal route was calculated by air traffic simulator SAAM by Eurocontrol:
 - -19 % less climate impact
 - 1 % longer flight time
 - 14 % more fuel
 - 22 % more NO_x
 - 10 % more costs



Frömming, et al. 2013 Grewe, et al. 2014, 2017 Matthes, et al. 2012,2017

Research objectives

- To investigate an eco-efficient aircraft routing strategy that reduces the climate impact of global air traffic over the next few decades
- To estimate its mitigation gain for different aircraft routing strategies

Methodologies

- Chemistry-climate model EMAC (ECHAM5/MESSy 2.54)
 Roeckner et al., 2006
 Jöckel 2010, 2016
- Submodel AirTraf 2.0 Yamashita, Kern et al. 2020
 - 9 routing strategies (called options)
 - Trajectory optimization (3D)
 - Geographic location, altitude, time of released non-CO₂ emissions/contrails are considered
 - Simplifications:
 - Only cruise flight phase
 - No potential conflicts of flight trajectories
 - No operational constraints from ATC

Aircraft routing options

- 0 Great circle
- 1 Flight time
- 2 Fuel use
- 3 NO_x emission
- 4 H₂O emission
- 5 Contrail formations
- 6 Simple operating cost
- 7 Cash operating cost
- 8 Climate impact



Flight trajectory optimization

- A trajectory (candidate) is created by B-spline curve with 11 design variables: 6 (geographical location), 5 (altitude)
- · Waypoints are automatically generated
- GA evaluates single objective function and finds out one optimal trajectory to minimize objective function value
 Objective function value





1-day air traffic simulations over the North Atlantic

Routing options	Great circle	Cost/Climate/Others options	
ECHAM5 Resolution	T42/L31ECMWF (2.8° × 2.8°)		
Duration / Time step	Dec.01.2015 - Dec.02.2015 / 12 min		
Waypoints	101		
Flight altitude change	Fixed FL350	FL290 - FL410	
Flight plan	103 Transatlantic flights by REACT4C Project		
	(Eastbound 52/Westbound 51)		
Aircraft / Engine type	A330-301 / CF6-80-E1-A2 (2GE051)		
EIH ₂ O [g(H ₂ O)/kg(fuel)]	1,230 (IPCC 1999)		
Load factor	0.62 (ICAO 2009)		
Fuel price [USD/USG]	1.545 (IATA 2017)		
Unit time cost [USD/h]	2710.0 (Boeing 2015)		
Mach number	0.82	0.82 (A330-301, Eurocontrol 2011)	
Optimization	-	Min. f (single-objective optimization)	
Design variable	-	11 (Location 6/Altitude 5)	
Generation number		100	
Population size	_	100	





Flight characteristics Dec. 1 2015, 103 North Atlantic flights (A330-301)

- Trade-off exists between operational cost and climate impact
- Climate-optimized routing can reduce expected climate impact (ATR20), compared to cost-optimized routing
 - Climate option: -67.9 % ATR20, +9.8 % COC → 0.13 [US Mil\$/10⁻⁷K]









Summary – research topics for further collaborations

- 1. To detect some unique points of the nondominated solutions and visualize the structure of nondominated fronts
- 2. To examine how much nondominated fronts vary under different weather conditions
- 3. To develop a decision-making method in EMAC/AirTraf

Elimination of B₂ singularities

Takahiro YAMAMOTO

Department of Mathematics, **T**okyo **G**akugei **U**niversity

8 January, 2021

Fiber Topology Meets Applications Kyushu University/ ZOOM

— Elimination of bdry singularities — 0/31



Joint Works with Computer Sci.:

(By O. Saeki, S. Takahashi, D. Sakurai, H. Wu, K. Kikuchi, H. Carr, D. Duke, **Y**) (1) **Visualizing Multivariate Data Using Singularity Theory**, The Impact of Applications on Mathematics, Proceedings of Forum of Mathematics for Industry 2013, 51–65.

(2) Interactive Visualization for Singular Fibers of Functions $f: \mathbb{R}^3 \rightarrow \mathbb{R}^2$, IEEE TRANSACTIONS ON VISUALIZATION AND COMPUTER GRAPHICS, VOL. 22, NO. 1, 945–954, JANUARY 2016.

! Contributed by classifying singular fibers of stable maps $M^3 \to \mathbb{R}^2$ of cpt 3-mfds with boundary...

(D. Sakurai, Y) (3) Visually Evaluating the Topological Equivalence of Bounded Bivariate Fields, accepted for publication in the Springer book on Topological Data Analysis following the TopoInVis 2019 workshop.

! Contributed by constructing the theory and calculating example by using "jc-net".

— Elimination of bdry singularities — 2/31

§ 0-2 Introduction

In this talk, mfd.s and maps are smooth of class C^∞ unless otherwise stated.

A smooth stable map, stable map for short is a smooth map s.t. any maps close to the map are equivalent to the map, $f: M \to P$ of a cpt mfd M with bdry into a surface P without bdry admits singularities which consists of

folds, cusps, bdry folds, B_2 (and bdry cusps if dim $M \ge 3$).

In this talk, we

 $\begin{cases} \textbf{give a set of local moves for stable maps } f \colon M \to P, \text{ and} \\ \textbf{study the question: Given } M, \text{ a stable map } f \colon M \to \mathbb{R}^2, \\ \text{can } f \text{ be deformed to a stable map } g \text{ s.t. } g \text{ has } \textbf{NO} B_2 \text{ singularities.} \end{cases}$

! It is a fundamental problem if generic singularities are eliminated by generic homotopy in singularity theory of smooth maps.

— Elimination of bdry singularities — 3/31

Contents:

- § 1. Preparation
- § 2. Local moves for $M \to P$
- \S 3. Main-Theorem and Outline of a proof of Main-Theorem

How to eliminate B_2 singularities (The case of dim M = 2):







Here are pictures of a cup taken from different angles. The picture on the left is taken from an oblique angle, and the object has **TWO** B_2 points. The pictures on the right shows the object viewed from directly above, and the picture in the middle is at in-between angle of the pictures on both sides. In the middle picture, the pair of B_2 points comes in collision, and the object eventually has **NO** B_2 points in the left picture.

- Elimination of bdry singularities - 4/31







Example: stable maps of cpt surf.s with bdry into the plane







 In order to study Reeb spaces of a stable map $f: M \to P$ of a 3-mfd with bdry into a surf., we recall the the behaviour of level set near a singular point $p \in S(f) \cup S(f|_{\partial})$:





§ 2 Local moves

Consider a C^{∞} homotopy τ as a path τ : $[0,1] \to C^{\infty}(M, P^2)$, $\tau(t) = \tau_t$.

By using the parametrized transversality theorem, we call a C^{∞} homotopy τ a **generic homotopy** if it satisfies the following condition: $\exists t_1, \ldots, t_k \in (0, 1)$ such that

(1) $\forall t \in (0,1) \setminus \{t_1, \dots, t_k\}, \tau_t \text{ is a stable map,}$ (2) $\forall t_i, (i = 1, \dots, k), \exists_1 y_i \in P^2 \text{ and}$ $\exists S_i \subset \tau_{t_i}^{-1}(y)$: a finite set, s.t. $\tau_{t_i} \colon (N, S_i) \rightarrow (P^2, y_i)$ is equivalent to one of the germs below: I For $t_{i-1} < t_i < t_i' < t_{i+1}$, the deformation of τ_t into τ_t' is given by the local moves given below.

- Elimination of bdry singularities - 14/31













§ 3 Main-Theorem and Outline of a proof of Main-Theorem: Let $H^n = \{(x_1, ..., x_n) \in \mathbb{R}^n \mid x_n \ge 0\}$. Then, $\partial H^n = \{x_n = 0\}$. Lemma 1.4(Characterization of B_2 Singularity) $f: (H^{(n)}, 0) \to (\mathbb{R}^2, 0)$ a smooth map germ. $f: (H^{(n)}, 0) \to (\mathbb{R}^2, 0)$ a smooth map germ. f: (n=2) f is B_2 singularity $\stackrel{\text{iff}}{\Leftrightarrow} \begin{cases} f \text{ is fold if we ignore the bdry,} \\ f|_\partial \text{ is immersion,} \\ S(f) \pitchfork \partial \text{ at } 0. \end{cases}$ f: (n=3) f is B_2 singularity $\stackrel{\text{iff}}{\Leftrightarrow} \begin{cases} f \text{ is fold if we ignore the bdry,} \\ f|_\partial \text{ is fold,} \\ S(f) \pitchfork \partial \text{ at } 0. \end{cases}$ I For a stable map $f: M \to P$ of an n-mfd (n = 2, 3) with bdry into a surf. B_2 singularities are initial and terminal points of an arc component of S(f). It implies

Lemma 1.5 M: a cpt *n*-mfd (n = 2, 3) with bdry, *P*: a surf. without bdry $f: M \to P$: a stable map, $S(f|_{\partial}) = S_1 \cup \cdots \cup S_{\ell}, \ \partial M = \partial_1 \cup \cdots \cup \partial_k$ the decomposition of $S(f|_{\partial}), \ \partial M$ into connected components resp. Then, we have the items below: If dimM = 2, then $\#\{B_2\}$ is even on $\forall \partial_j$. If dimM = 3, then $(1) \ \#\{B_2\} + \#\{$ bdry cusps $\}$ is even on $\forall S_i, \ \beta_i$ $(2) \ \#\{B_2\}$ and $\chi(\partial_j)$ have the same parity on $\forall \partial_j$. β_i It yields that if dimM = 3 and $\exists \partial_j \subset \partial M$ s.t. $\chi(\partial_j)$ is odd, then any stable map $f: M \to P$ has at least ONE B_2 point on $\partial_j \subset \partial N$. For the cases dim $M \ge 4$, there are topological constraints of # of B_2 as well.

- Elimination of bdry singularities - 22/31

Main-Theorem A M: a cpt n-mfd (n = 2, 3) with bdry, $\partial M = \partial_1 \cup \cdots \cup \partial_k$, f: $M \rightarrow P$: a stable map. (n = 2): f is homotopic to a stable map $g: N \rightarrow P^2$ having NO B_2 points, namely, f is homotopic to a stable map $g: N \rightarrow P^2$ which is SUBMERSION on a sufficiently small nbd of ∂ . (n = 3): Then, f is homotopic to a stable map $g: N \rightarrow P^2$ s.t. $\left\{ \begin{array}{l} \text{if } \chi(\partial_j) \stackrel{(2)}{=} 0, \text{ then } g \text{ has NO } B_2 \text{ points on } \partial_j, \\ \text{if } \chi(\partial_j) \stackrel{(2)}{=} 0, \text{ then } g \text{ has ONE } B_2 \text{ point on } \partial_j. \end{array} \right.$ In particular, if $\chi(\partial_j) \stackrel{(2)}{=} 0$ for $\forall \partial_j$, then f is homotopic to a stable map $g: M \rightarrow P^2$ which is SUBMERSION on a sufficiently small nbd of ∂M . I The Theorem implies that g is obtained by applying local moves above finitely from f. I We had obtained similar type theorem for the cases dim $M \ge 4$.

— Elimination of bdry singularities — 23/31

To obtain a stable map $g: M \to P$ having NO B_2 pt.s, f is deformed by local moves given above following four steps below.



! def/indef- C_3 -homotopy replace an def/indef- B_2 point with a bdry cusp point.

- Elimination of bdry singularities - 25/31



$$\begin{split} & \begin{array}{l} f_{3} \sim_{homotopic} f_{4}, \\ \text{where } f_{4} \text{ is a stable map s.t.} \\ & \text{if } \chi(\partial_{j}) \stackrel{(2)}{=} 0, \text{ then } \#\{B_{2}\} \text{ points is ZERO or TWO on } \forall S_{i} \subset S(f_{4}|_{\partial_{j}}), \\ & \text{if } \chi(\partial_{j}) \stackrel{(2)}{\neq} 0, \text{ then } \exists_{1}S_{i'} \subset S(f_{4}|_{\partial_{j}}), \\ & \text{s.t. } \#\{B_{2}\} \text{ points is ONE on } S_{i'}, \\ & \#\{B_{2}\} \text{ points is ZERO or TWO on } \forall S_{i} \subset S(f_{4}|_{\partial_{j}}) \setminus S_{i'}. \\ \\ & \text{Then, by applying Steps 2 and 3 to } f_{4}, \text{ we obtain the desired stable map } g: M \rightarrow P^{2}. \\ \\ & \text{To apply STEP 4, there are two cases for } f_{3}: \\ & - \text{Elimination of bdry singularities} - 27/31 \end{split}$$

Case (1) $\gamma \subset S(f)$: an arc whose bdry is a pair of B_2 points $p_1, p_2 \in \partial_j$ s.t. $p_1 \in S_{j_1}$ and $p_2 \in S_{j_2}$ $(S_{j_1}, S_{j_2} \subset S(f|_{\partial_i}))$.

Assume that for p_1 and p_2 , \exists bdry cusps in the sufficiently small nbd of p_1 , p_2 resp. Then, by using C_3 -homotopy if necessary, we obtain the following lemma.

Let η be an embedded curve transversal to $S(f_3|_{\partial})$ with $\eta(0) = p_{\lambda}$, $\eta(1) = p_{\mu}$ s.t. there are NO bdry cusps f_3 on $\eta((0,1))$. Then, \exists a homotopy s.t. the homotopy does NOT change f_3 on $N \setminus N(\eta)$ and it change f_3 to a stable map which has no cusps on $N(\eta)$.

- Elimination of bdry singularities - 28/31

Case (2): $\gamma_{\lambda} \subset S(f)$: arc whose bdry are $p_{\lambda_1} \in \partial_{j_1}$ and $p_{\lambda_2} \in \partial_{j_2}$, $\gamma_{\mu} \subset S(f)$: arc whose bdry are $p_{\mu_1} \in \partial_{j_1}$ and $p_{\mu_3} \in \partial_{j_3}$.

Assume that for p_{λ_1} and $p_{\lambda_2}, p_{\mu_1}, p_{\mu_3}$, there are bdry cusps in the sufficiently small nbhd of p_{λ_1} and $p_{\lambda_2}, p_{\mu_1}, p_{\mu_3}$ resps. Then, by using C_3 -homotopy and Sw.-homotopy if necessary, we obtain the following Lemma.

Let η be an embedded curve transversal to $S(f_3|_{\partial})$ with $\eta(0) = p_{\lambda_1}$, $\eta(1) = p_{\mu_1}$ s.t. there are NO bdry cusps on $\eta((0,1))$. Then, \exists a homotopy s.t. the homotopy does NOT change f_3 on $N \setminus N(\eta)$ and it change f_3 to a stable map which has no cusps on $N(\eta)$.

— Elimination of bdry singularities — 29/31









Lefschetz Pencil

§1. Broken Lefschetz Fibrations §2. Indefinite Fibrations §3. Moves §4. Constructions §5. Trisections

Lefschetz pencil (LP) is a LF $M^4 \setminus B \to S^2$ for a finite set $B \neq \emptyset$, where it has complex local model

$$(z_1, z_2) \mapsto z_1/z_2 \in \mathbf{C} \cup \{\infty\} = S^2.$$

(Blowing-up the points in B, we get a LF $M^4 \sharp (\sharp^{|B|} \overline{\mathbb{C}P^2}) \to S^2$. Conversely, blowing-down (-1)-sections for a LF, we get a LP.)

Donaldson, Gompf, Late 90's ∃Lefschetz pencil ⇔ ∃symplectic structure

Symplectic structure: $\omega \in \Omega^2(M^4)$, $d\omega = 0$, non-degenerate ($\omega^2 > 0$)

∧ Symplectic structures are "**HARD**" to find.



§1. Broken Lefschetz Fibrations §2. Indefinite Fibrations §3. Moves §4. Constructions §5. Trisections

\S **2. Indefinite Fibrations**

Indefinite Fibration

§1. Broken Lefschetz Fibrations §2. Indefinite Fibrations §3. Moves §4. Constructions §5. Trisections

A singularity of $M^4 \to \Sigma^2$ is a fold if it is locally given by $(t, x_1, x_2, x_3) \mapsto (t, \pm x_1^2 \pm x_2^2 \pm x_3^2)$. A singularity is a **cusp** if it is locally given by $(t, x_1, x_2, x_3) \mapsto (t, x_1^3 + tx_1 \pm x_2^2 \pm x_3^2)$.

Whitney, Thom, ~1950's Any smooth map can be approximated by a map with only fold and cusp singularities. (Such $f: M^4 \to \Sigma^2$ is called a generic map.)

f is an **indefinite generic map** if its folds and cusps are all **indefinite**. $f: M^4 \rightarrow \Sigma^2$ is an **indefinite fibration** (**IF**) if it has at most indefinite folds, indefinite cusps, and **Lefschetz** singularities.





	\$1. Broken Lefschetz Fibrations $$2.$ Indefinite Fibrations $$3.$ Moves $$4.$ Constructions $$5.$ Trisections	
	§3. Moves	
е.		•





Moves for Indefinite Fibrations

§1. Broken Lefschetz Fibrations §2. Indefinite Fibrations §3. Moves §4. Constructions §5. Trisections

Definition 3.1 A move for an IF $f: M^4 \to \Sigma^2$ is a smooth 1-parameter family $f_t: M^4 \to \Sigma^2$, $t \in [0, 1]$, of "mostly" IFs, with $f_0 = f$, which modifies the base diagram only in a nbhd of a point in Σ^2 . (Except for finitely many t's, f_t is an IF.)

Definition 3.2 A move is **always-realizable** if, given a local configuration of a base diagram, we can **always** find a 1-parameter family as above that realizes the relevant base diagram change.












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An Efficient Triangulation for Extruded Spatiotemporal Prism Meshes

<u>Akito Fujii</u> Kenji Ono Daisuke Sakurai

1

2

Triangulation of spatiotemporal meshes

- Contour tree requires spatial triangulation
- Thanks to in-situ analysis, very fine temporal resolution!
- We want to triangulate spatiotemporal meshes without breaking pre-defined spatial triangulation



Our proposed spatiotemporal triangulation!

- Implicit :
 - On-the-fly from the information in the node
 - Ex. crystalline triangulation
- Features :
 - Arbitrary dimension
 - Arbitrary spatial triangulation
 - Extrude to time direction
 - Fast
 - Parallelism
- Assume temporally static meshes



4

Conclusion

- We propose triangulation algorithm for extruded prism meshes
 - This algorithm can apply meshes in arbitrary dimensions
- Low cost
 - Cost does not depend on the number of vertices
 - We do not have to store triangulation of meshes
 - Look up table for the triangulation only depends on the number of dimension

6

• Easy to implement

「マス・フォア・インダストリ研究」シリーズ刊行にあたり

本シリーズは、平成23年4月に設立された九州大学マス・フォア・インダストリ研究所 (IMI)が、平成25年4月に共同利用・共同研究拠点「産業数学の先進的・基礎的共同研究 拠点」として、文部科学大臣より認定を受けたことにともない刊行するものである.本シ リーズでは、主として、マス・フォア・インダストリに関する研究集会の会議録、共同研 究の成果報告等を出版する.各巻はマス・フォア・インダストリの最新の研究成果に加え、 その新たな視点からのサーベイ及びレビューなども収録し、マス・フォア・インダストリ の展開に資するものとする.

> 平成 30 年 10 月 マス・フォア・インダストリ研究所 所長 佐伯 修

Fiber Topology Meets Applications

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